

ARP and RARP

Logical Addresses

- The hosts and routers are recognized at the network level by their *logical addresses*
 - A **logical address** is an internet address
 - Called a *logical* address because it is usually implemented in software
 - The logical addresses in the TCP/IP are called **IP address** and are 32 bits long

Physical Address

- However, hosts/routers are recognized at the physical layer by their *physical address*
 - A **physical address** is an local address

- Called a *physical* address because it is usually implemented in hardware

- Examples
 - 48-bit MAC addresses in Ethernet

Translation

- We need both the physical address and the logical address for packet delivery
- Thus, we need to be able to map a logical address to its corresponding physical address and vice versa
- Solutions
 - *Static mapping*
 - *Dynamic mapping*

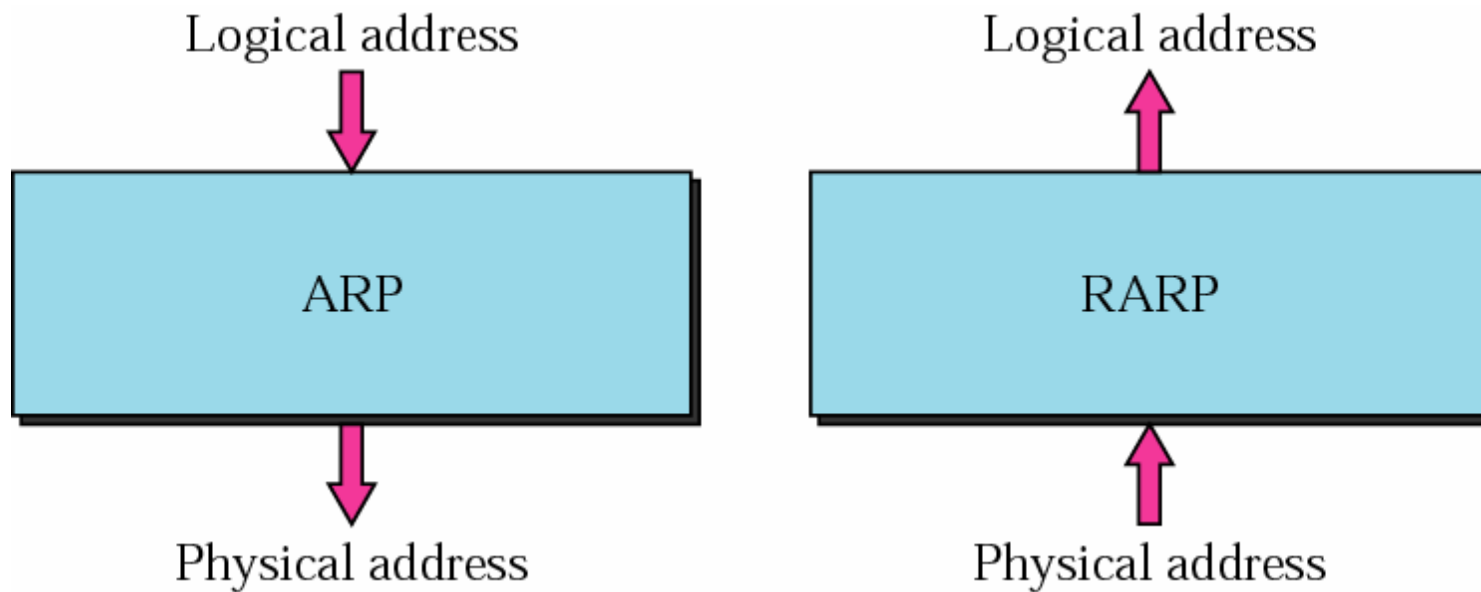
Static Mapping

- Create a table that associates a logical address with a physical address and store in each machine
- However, physical addresses may change
 - A machine could change its NIC resulting in a new physical address
 - In some LANs, such as LocalTalk, the physical address changes every time the computer is turned on
 - A mobile station can move from one physical network to another, resulting in a change in its physical address

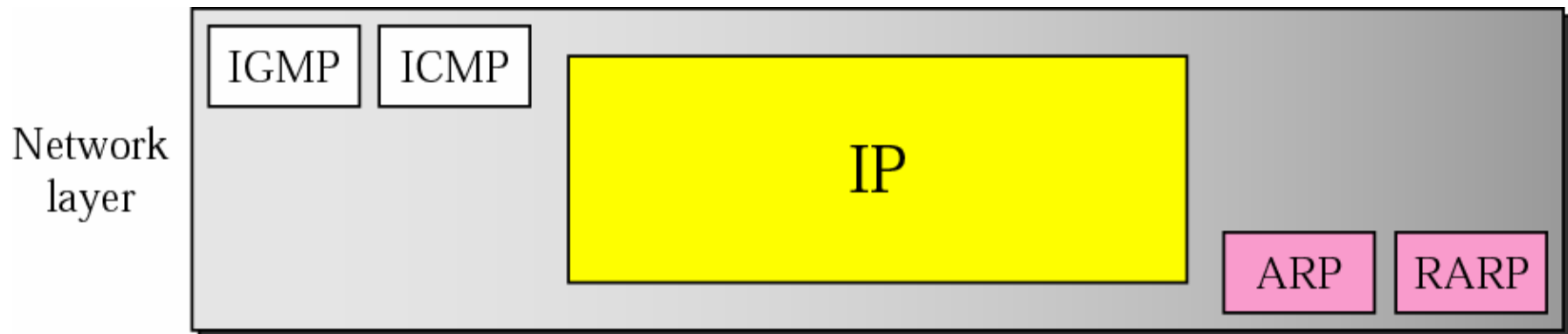
Dynamic Mapping

- Use a protocol to find another address
- ARP: Address Resolution Protocol
 - Map a logical address to a physical address
- RARP: Reverse Address Resolution Protocol
 - Map a physical address to a logical address

ARP and RARP



Position of ARP and RARP in TCP/IP Protocol Suite



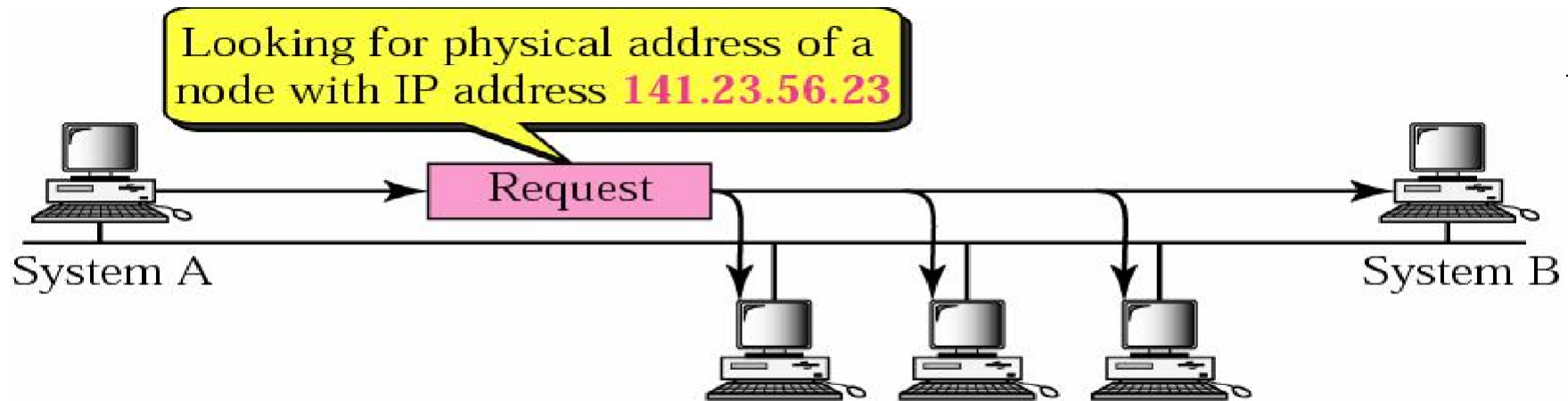
ARP Operation

- To find the physical address of another host or router on its network
 - Send an ARP request message
- ARP request message
 - The physical address of the sender
 - The IP address of the sender
 - The physical address of the receiver is *0s*
 - The IP address of the receiver

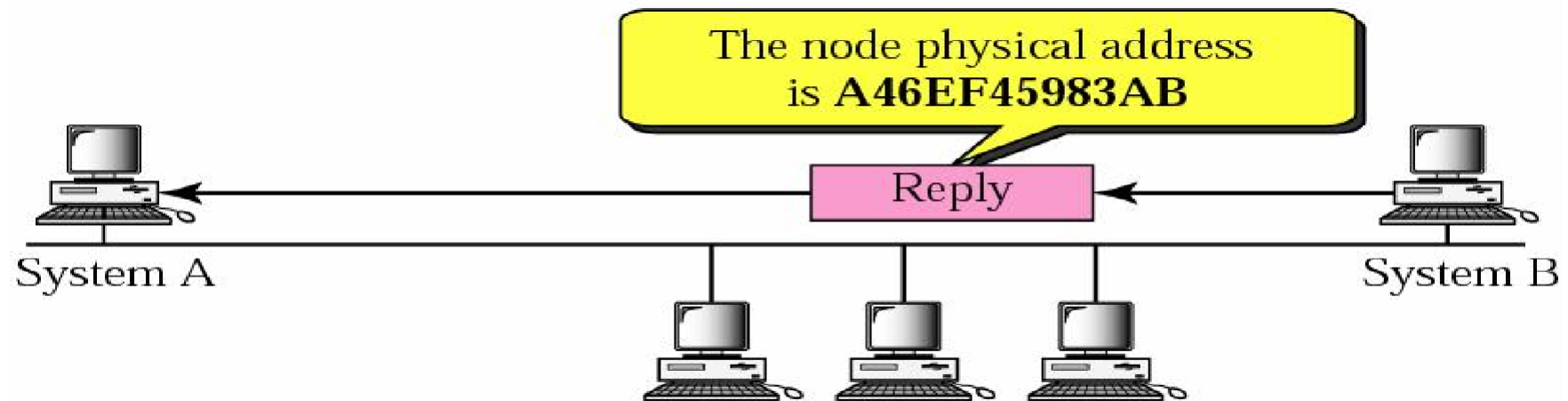
ARP Operation (Cont.)

- Then, ARP request message is broadcast by the physical layer
 - For example: in Ethernet, MAC header's destination address is all *1s* (broadcast address)
 - Received by every station on the physical network
- The intended recipient send back an ARP reply message
 - ARP reply message packet is *unicast*

ARP Operation



a. ARP request is broadcast



b. ARP reply is unicast

ARP Packet

Hardware Type		Protocol Type
Hardware length	Protocol length	Operation Request 1, Reply 2
Sender hardware address (For example, 6 bytes for Ethernet)		
Sender protocol address (For example, 4 bytes for IP)		
Target hardware address (For example, 6 bytes for Ethernet) (It is not filled in a request)		
Target protocol address (For example, 4 bytes for IP)		

Packet Format

- HTYPE (Hardware type)
 - 16-bit field defining the underlying type of the network
 - Ethernet is given the type 1
 - ARP can be used on any physical network
- PTYPE (Protocol type)
 - 16-bit field defining the protocol
 - IPv4 is 0800_{16}
 - ARP can be used with any higher-level protocol

Packet Format (Cont.)

- HLEN (Hardware length)
 - 8-bit field defining the length of the physical address in bytes
 - Ethernet has the value of 6
- PLEN (Protocol length)
 - 8-bit field defining the length of the logical address in bytes
 - IPv4 has the value of 4
- OPER (Operation)
 - 16-bit field defining the type of packet
 - (1) = ARP request, (2) = ARP reply

Packet Format (Cont.)

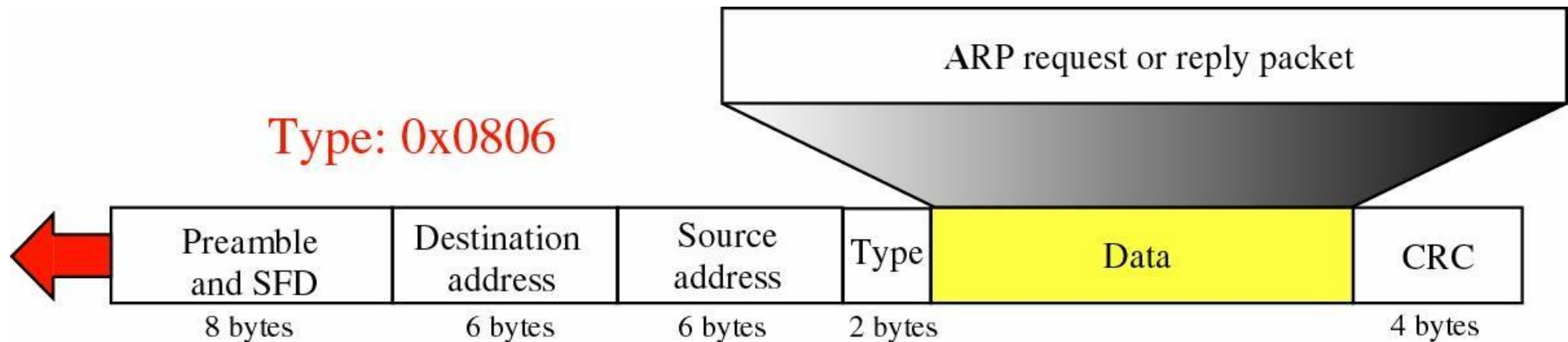
- SHA (Sender hardware address)
 - A variable-length field defining the physical address of the sender

- SPA (Sender protocol address)
 - A variable-length field defining the logical address of the sender

Packet Format (Cont.)

- **THA (Target hardware address)**
 - A variable-length field defining the physical address of the target
 - For an ARP request operation packet
 - This field is all 0s
- **TPA (Target protocol address)**
 - A variable-length field defining the logical address of the target

Encapsulation of ARP Packet



- An ARP packet is encapsulated directly into a data link frame
- Type field indicates that the data carried by the frame is an ARP packet

Operations

- The sender knows the target's IP address
- IP asks ARP to create an ARP request message
 - The sender physical address
 - The sender IP address
 - The target physical address field is filled with 0s
 - The target IP address
- The message is passed to the data link layer to encapsulate in a data link frame
 - Physical destination address is broadcast address

Operations (Cont.)

- Every host or routers receives the frame and since the destination address is broadcast, pass it to the ARP
 - All machines' ARP except the one targeted drop the packet
- The target reply with an ARP reply message that contains its physical address and is unicast
- The sender receives the reply message and knows the target's physical address

Four Cases to Use ARP

- **Case 1:** The sender is a host and wants to send a packet to another host on the same network
 - Use ARP to find another host's physical address
- **Case 2:** The sender is a host and wants to send a packet to another host on another network
 - Sender looks at its routing table
 - Find the IP address of the next hop (router) for this destination
 - Use ARP to find the router's physical address

Four Cases Using ARP: Case 1

Target IP address:
Destination address in the IP datagram

Sender



Host



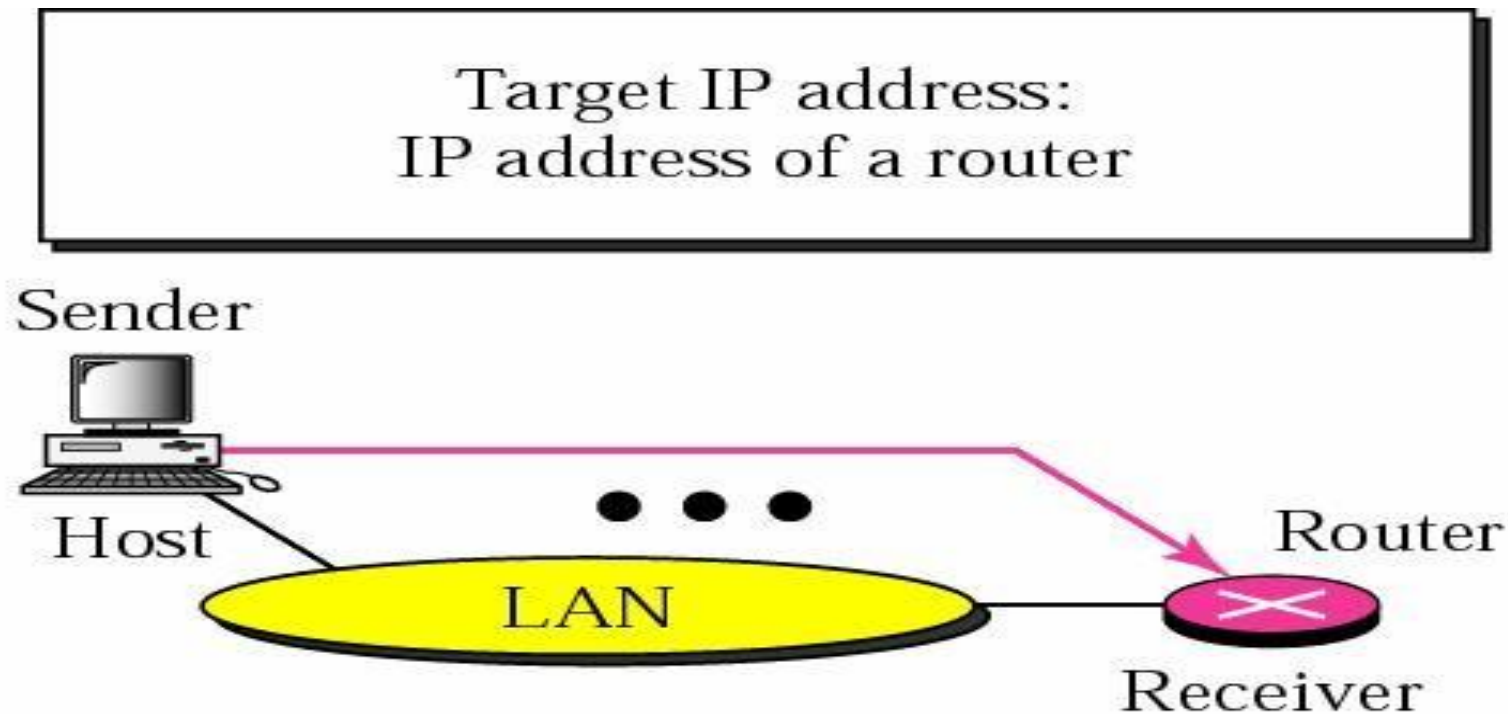
Host



Receiver

Case 1. A host has a packet to send to another host on the same network.

Four Cases Using ARP: Case 2

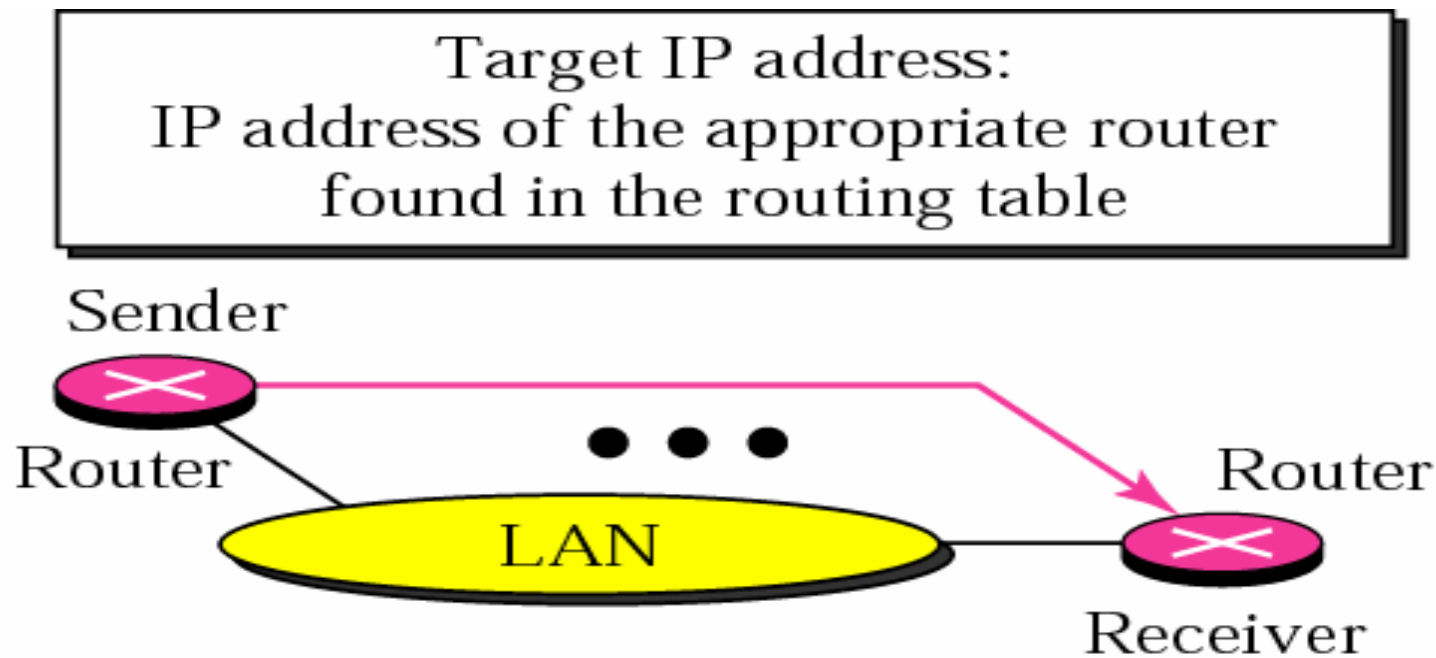


Case 2. A host wants to send a packet to another host on another network.
It must first be delivered to a router.

Four Cases to Use ARP (Cont.)

- **Case 3:** the sender is a router and received a datagram destined for a host on another network
 - Router check its routing table
 - Find the IP address of the next router
 - Use ARP to find the next router's physical address
- **Case 4:** the sender is a router that has received a datagram destined for a host in the same network
 - Use ARP to find this host's physical address

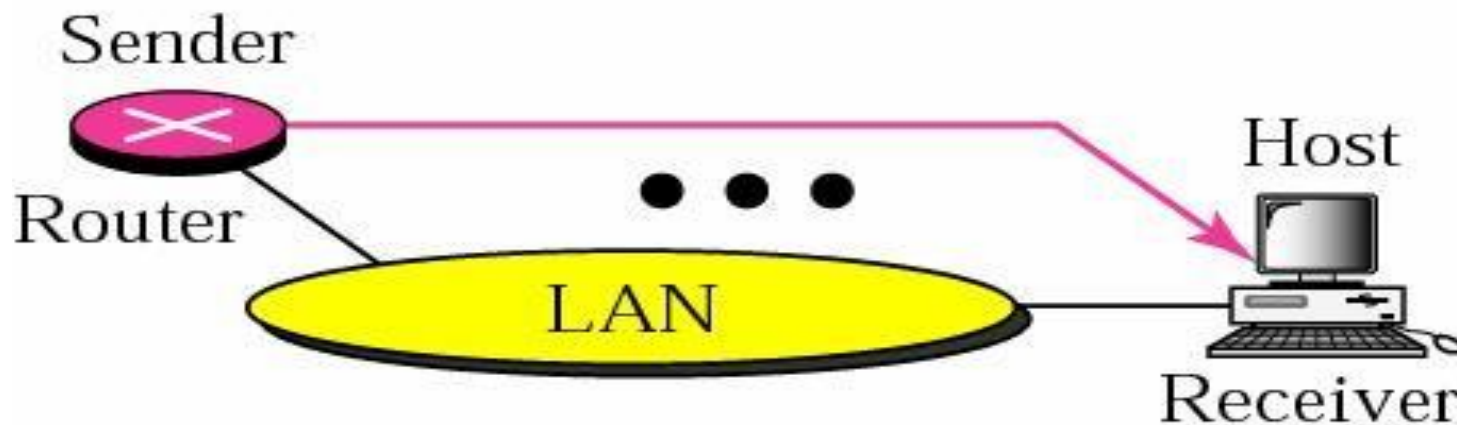
Four Cases Using ARP: Case 3



Case 3. A router receives a packet to be sent to a host on another network. It must first be delivered to the appropriate router.

Four Cases Using ARP: Case 4

Target IP address:
Destination address in the IP datagram



Case 4. A router receives a packet to be sent to a host on the same network.

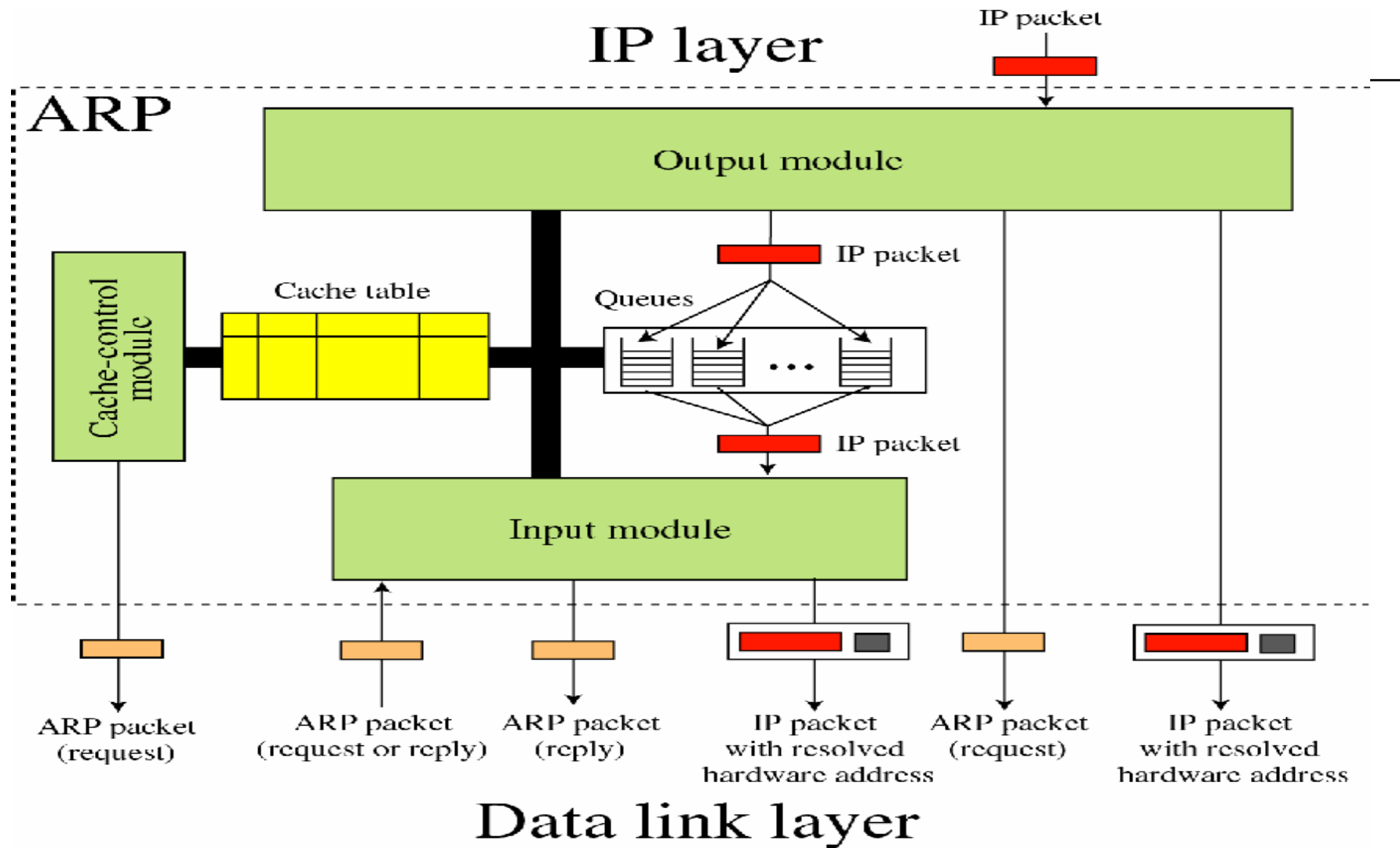
Proxy ARP

- Used to create a subnetting effect
- A router running a proxy ARP
 - Its ARP acts on behalf of a set of hosts
 - If it receives an ARP request message looking for the address of one of these host
 - The router sends an ARP reply announcing its own hardware (physical) address
 - After the router receives the actual IP packet
 - It sends the packet to the appropriate host or router

ARP Package

- Five components in an ARP package
 - A cache table
 - Queues
 - An output module
 - An input module
 - A cache-control module

ARP Components



Cache Table

- Inefficient to use ARP to each datagram destined for the same host or router
 - Introduce the cache table
- Cache table: an array of entries that contains the following's entries

Content of a Cache Table Entry

- State:
 - FREE: the time-to-live for this entry has expired
 - PENDING: a request for this entry has been sent, but the reply has not yet been received
 - RESOLVED: the entry is complete and valid
- Hardware type
- Protocol type
- Hardware length
- Protocol length
 - Above fields are all the same as in the ARP packet

Content of a Cache Table Entry (Cont.)

- Interface number
- Queue number: ARP uses numbered queues to enqueue the packet waiting for address resolution
- Attempts: the number of times an ARP request is sent out for this entry
- Time-out: the lifetime of an entry in seconds
- Hardware address: the destination hardware address
- Protocol address: the destination IP address

Queues

- ARP package maintains a set of queues to hold the IP packets while ARP tries to resolve the hardware address
- Packets for the same destination are usually enqueued in the same queue
- The output module sends unsolved packets into the queue
- The input module removes a packet from the queue and sends it, with the resolved physical address, to data link layer for transmission

Output Module

- Wait until an IP packet from the IP software
- Check the cache table if receiving a IP packet
 - If found and state = RESOLVED
 - Passed to the data link layer for transmission
 - If found and state = PENDING
 - Send packet to this queue and wait
 - If not found
 - Create an entry with state = PENDING
 - Create a queue and enqueue this packet
 - Send an ARP request

Input Module

- Wait until an ARP packet (request or reply) arrives and check the cache table
 - If found state = PENDING
 - Copy the target hardware address in the packet
 - Change the state to RESOLVED
 - Set the value of TIME-OUT for this entry
 - Dequeue the packets from the corresponding queue and set them to the data link layer

Input Module (Cont.)

- If found and state = RESOLVED
 - Copy the target hardware address in the packet
 - Set the value of TIME-OUT for this entry
 - This is because the target hardware address could have been changed
- If not found
 - Create a new entry and adds it to the table
- If the packet is a request
 - Send an ARP reply

Cache-Control Module

- Maintain the cache table by periodically check the cache table, entry by entry
- If state is PENDING
 - Increment the value of attempts by 1
 - If (attempts greater than maximum)
 - Change the state to FREE and Destroy the corresponding queue
 - Else
 - Send an ARP request

Cache-Control Module (Cont.)

- If state is RESOLVED
 - Decrement the value of time-out by the value of elapsed time
 - If (time-out \leq 0)
 - Change the state to FREE
 - Destroy the corresponding queue
- If state is FREE
 - Continue to the next entry

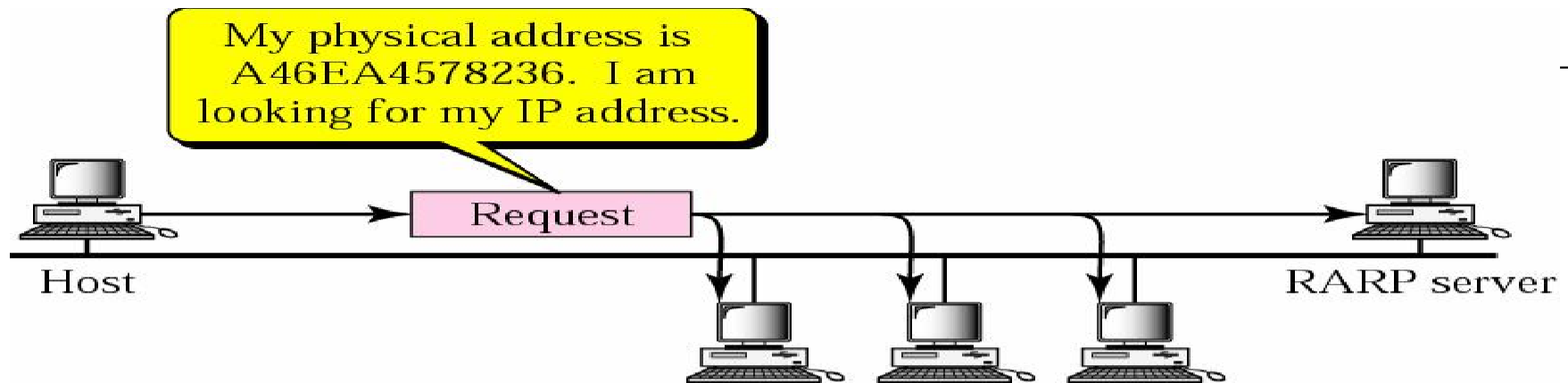
Original Cache Table

<i><u>State Queue Attempt Time-out Protocol Addr.</u></i>				<i><u>Hardware Addr.</u></i>	
R	5		900	180.3.6.1	ACAE32457342
P	2	2		129.34.4.8	
P	14	5		201.11.56.7	
R	8		450	114.5.7.89	457342ACAE32
P	12	1		220.55.5.7	
F					
R	9		60	19.1.7.82	4573E3242ACA
P	18	3		188.11.8.71	

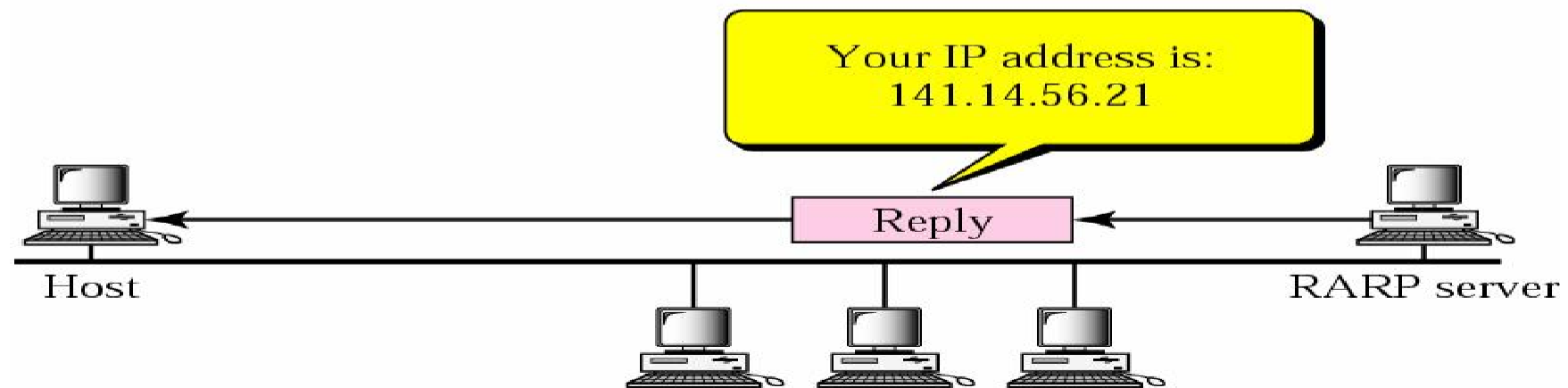
RARP

- A diskless machine is usually booted from ROM
- It cannot include the IP address
 - IP address are assigned by the network administrator
- Obtain its logical address by the physical address using the RARP protocol

RARP Operation



a. RARP request is broadcast



b. RARP reply is unicast

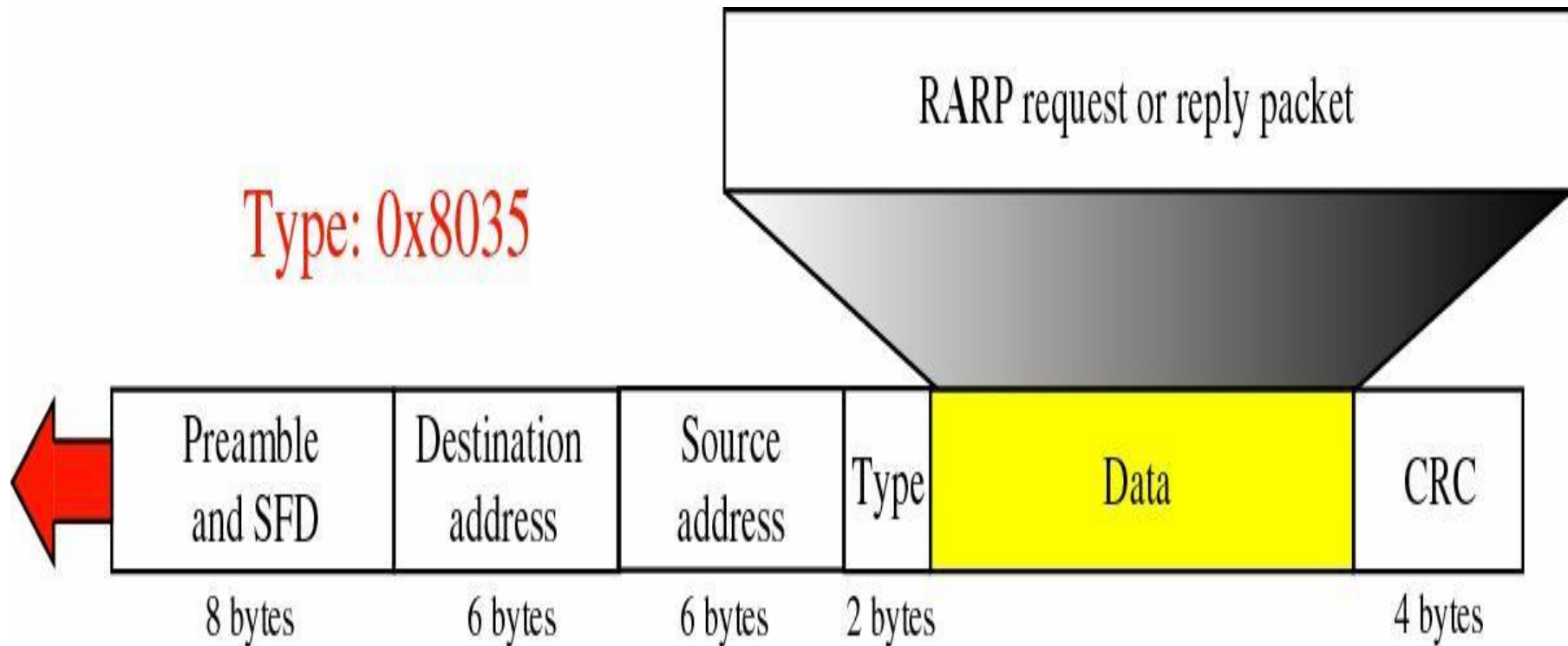
Packet Format

- The format of the RARP packet is the same as the ARP packet
- Except that the operation field is
 - Three for RARP request message
 - Four for RARP reply message

RARP Packet

Hardware type		Protocol type
Hardware length	Protocol length	Operation <i>Request 3, Reply 4</i>
Sender hardware address (For example, 6 bytes for Ethernet)		
Sender protocol address (For example, 4 bytes for IP) (It is not filled for request)		
Target hardware address (For example, 6 bytes for Ethernet) (It is not filled for request)		
Target protocol address (For example, 4 bytes for IP) (It is not filled for request)		

Encapsulation of RARP Packet



Alternative Solutions to RARP

- When a diskless computer is booted, it needs more information in addition to its IP address
 - The subnet mask
 - The IP address of a router
 - The IP address of a name server
- RARP cannot provide this extra information
- Two protocols, BOOTP and DHCP, can be used instead of RARP