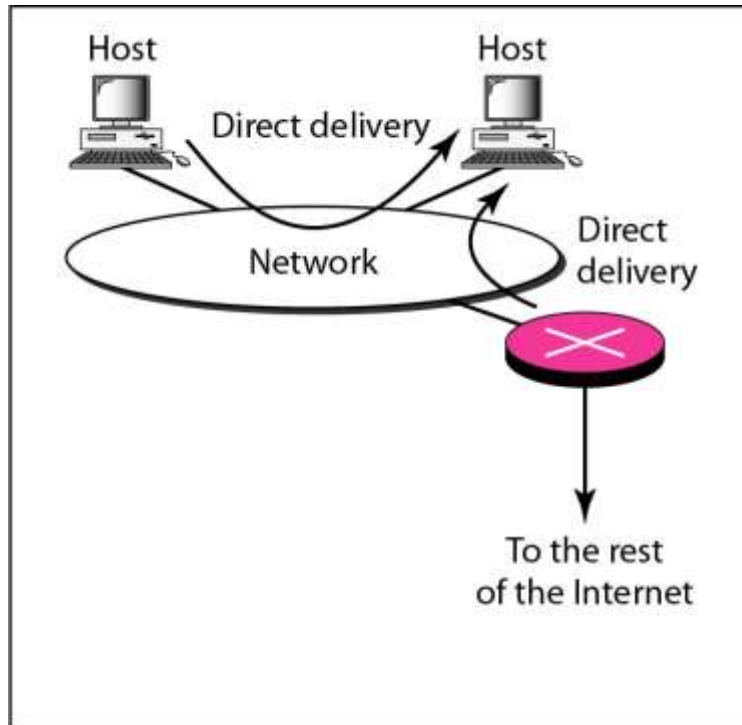

*Delivery, Forwarding in Network Layer
& Hierarchical Routing*

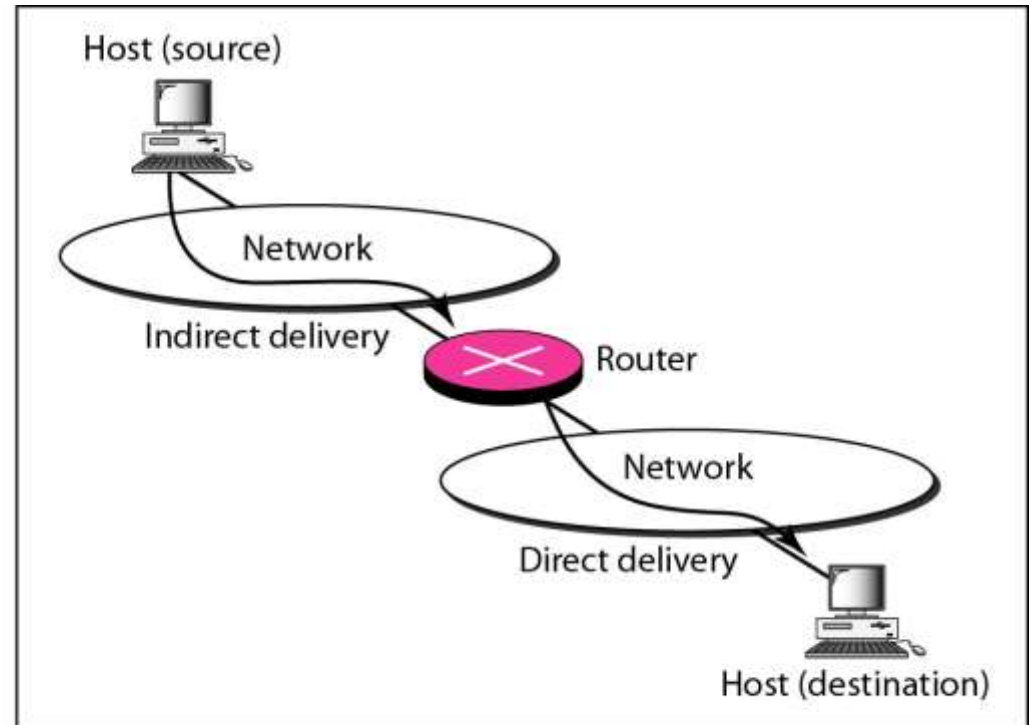
1. DELIVERY

The network layer supervises the handling of the packets by the underlying physical networks. We define this handling as the delivery of a packet.

Figure-1 *Direct and indirect delivery*



a. Direct delivery



b. Indirect and direct delivery

2. FORWARDING

Forwarding means to place the packet in its route to its destination. Forwarding requires a host or a router to have a routing table. When a host has a packet to send or when a router has received a packet to be forwarded, it looks at this table to find the route to the final destination.

Figure-2 *Route method versus next-hop method*

a. Routing tables based on route

| Destination | Route |
|-------------|----------------|
| Host B | R1, R2, host B |

Routing table
for host A

| Destination | Route |
|-------------|------------|
| Host B | R2, host B |

Routing table
for R1

| Destination | Route |
|-------------|--------|
| Host B | Host B |

Routing table
for R2

b. Routing tables based on next hop

| Destination | Next hop |
|-------------|----------|
| Host B | R1 |

| Destination | Next hop |
|-------------|----------|
| Host B | R2 |

| Destination | Next hop |
|-------------|----------|
| Host B | --- |

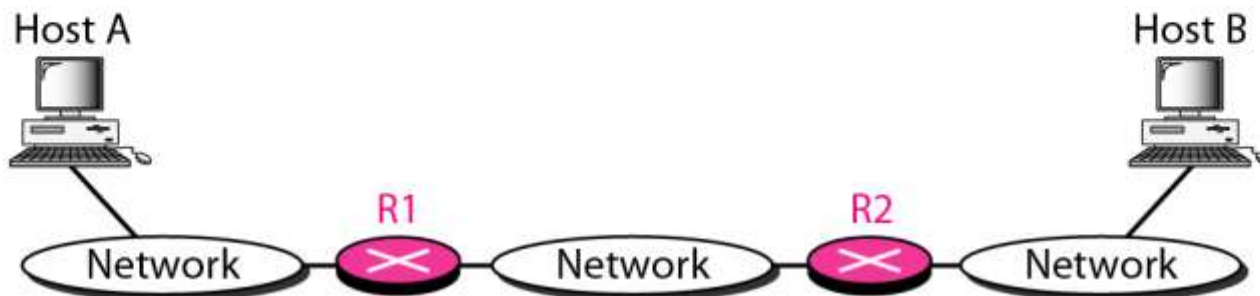


Figure-3 *Host-specific versus network-specific method*

Routing table for host S based on host-specific method

| Destination | Next hop |
|-------------|----------|
| A | R1 |
| B | R1 |
| C | R1 |
| D | R1 |

Routing table for host S based on network-specific method

| Destination | Next hop |
|-------------|----------|
| N2 | R1 |

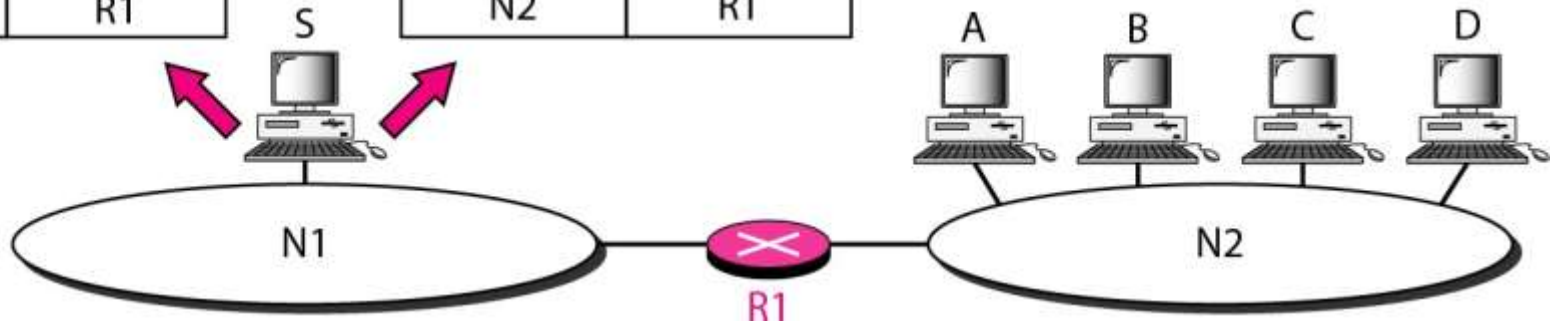


Figure-4 *Default method*

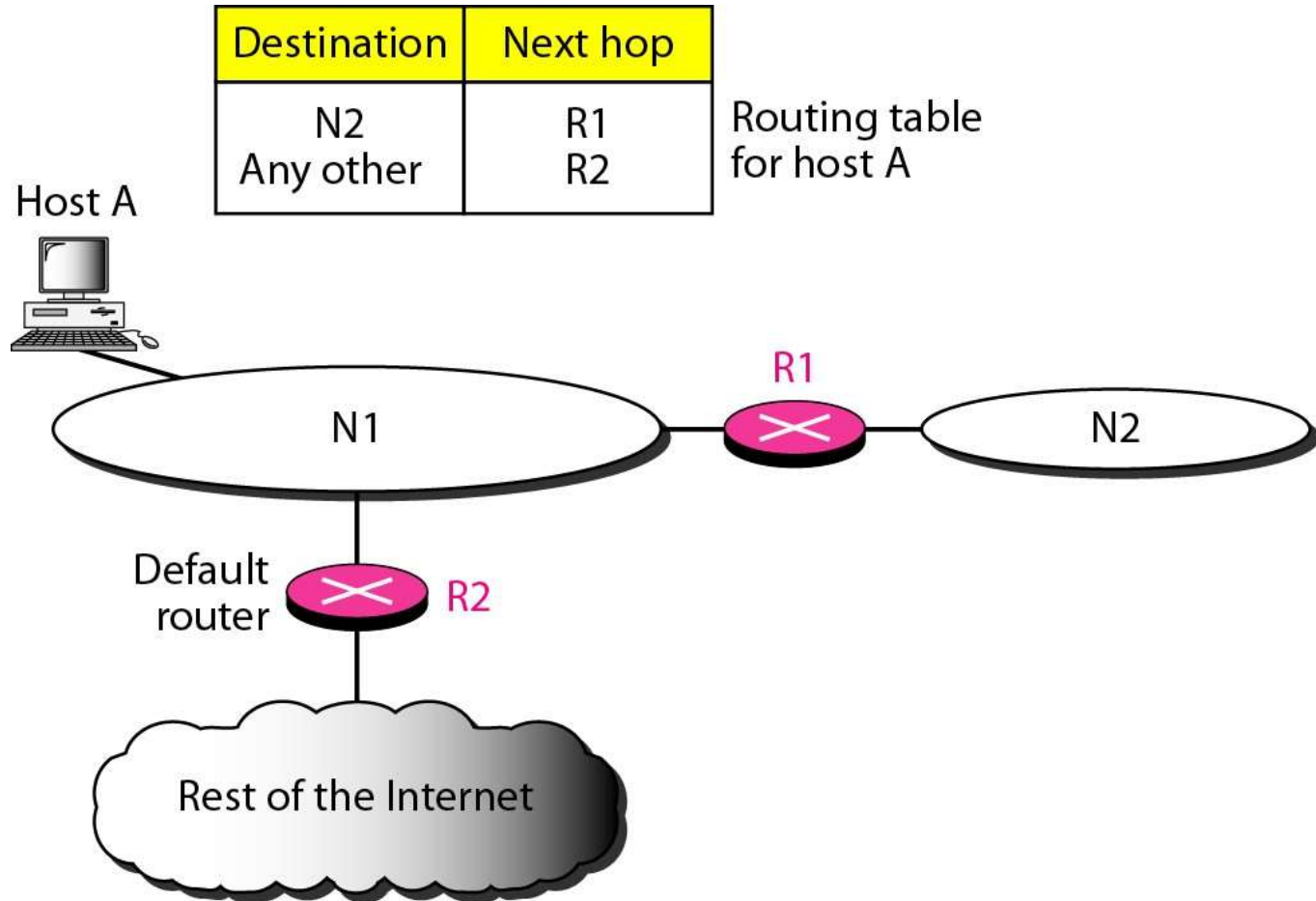
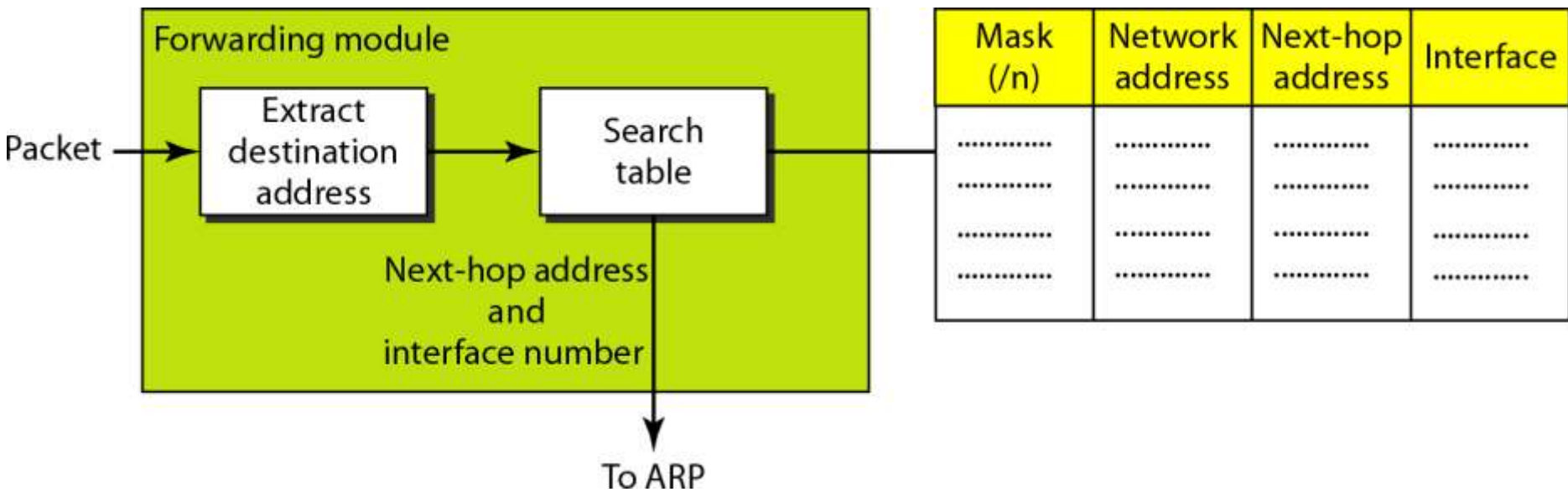


Figure-5

Simplified forwarding module in classless address





Note

In classless addressing, we need at least four columns in a routing table.



Example-1

Make a routing table for router R1, using the configuration in Figure-6.

Solution

Table-1 shows the corresponding table.

Figure-6 *Configuration for Example-1*

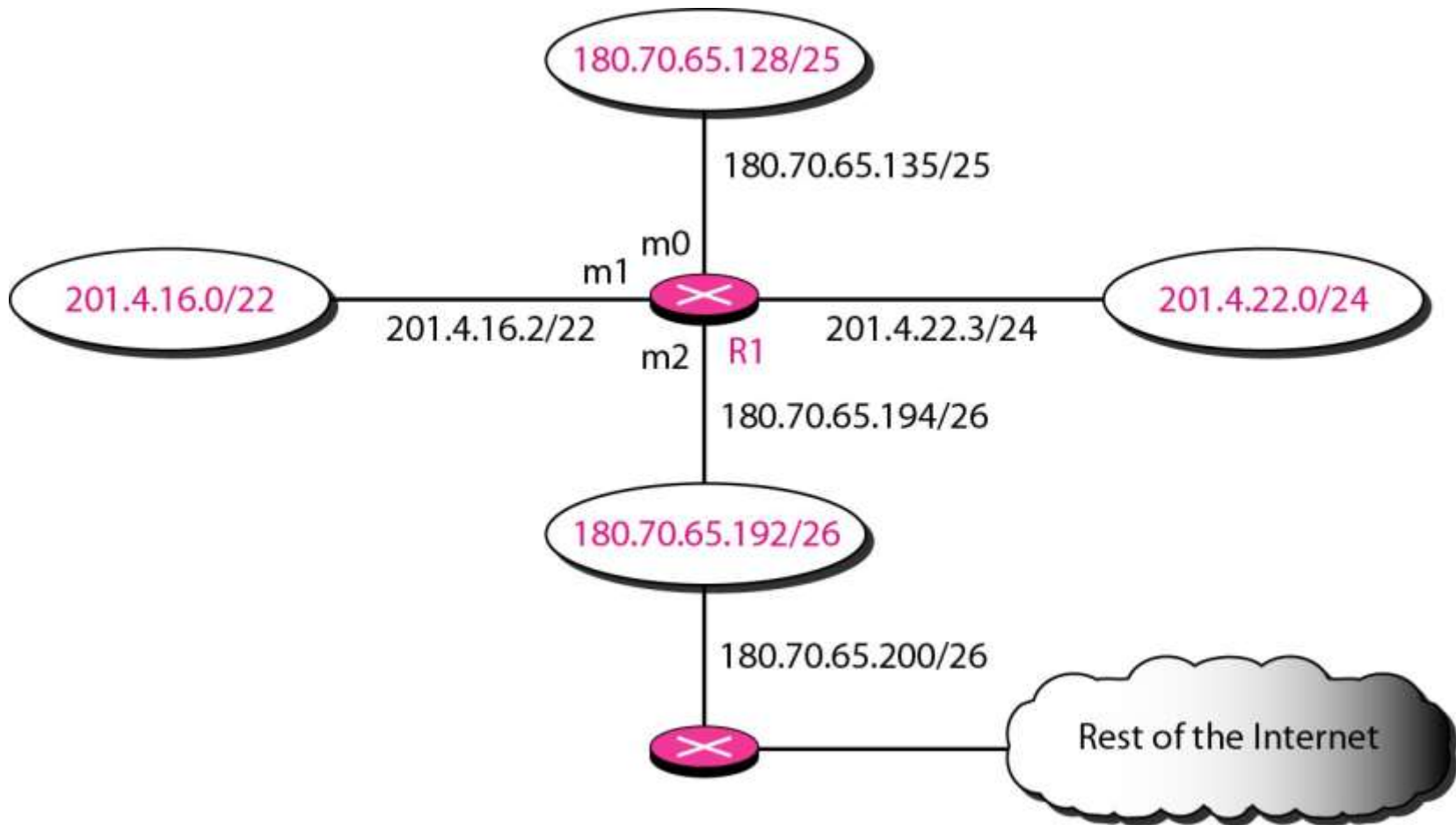


Table-1 *Routing table for router R1 in Figure-6*

| <i>Mask</i> | <i>Network Address</i> | <i>Next Hop</i> | <i>Interface</i> |
|-------------|------------------------|-----------------|------------------|
| <i>/26</i> | 180.70.65.192 | — | m2 |
| <i>/25</i> | 180.70.65.128 | — | m0 |
| <i>/24</i> | 201.4.22.0 | — | m3 |
| <i>/22</i> | 201.4.16.0 | | m1 |
| Any | Any | 180.70.65.200 | m2 |



Example-2

Show the forwarding process if a packet arrives at R1 in Figure-6 with the destination address 180.70.65.140.

Solution

The router performs the following steps:

- 1. The first mask (/26) is applied to the destination address. The result is 180.70.65.128, which does not match the corresponding network address.*
- 2. The second mask (/25) is applied to the destination address. The result is 180.70.65.128, which matches the corresponding network address. The next-hop address and the interface number m0 are passed to ARP for further processing.*



Example-3

Show the forwarding process if a packet arrives at R1 in Figure-6 with the destination address 201.4.22.35.

Solution

The router performs the following steps:

- 1. The first mask (/26) is applied to the destination address. The result is 201.4.22.0, which does not match the corresponding network address.*
- 2. The second mask (/25) is applied to the destination address. The result is 201.4.22.0, which does not match the corresponding network address (row 2).*

Example-3 (continued)

- 3. The third mask (/24) is applied to the destination address. The result is 201.4.22.0, which matches the corresponding network address. The destination address of the packet and the interface number m3 are passed to ARP.*



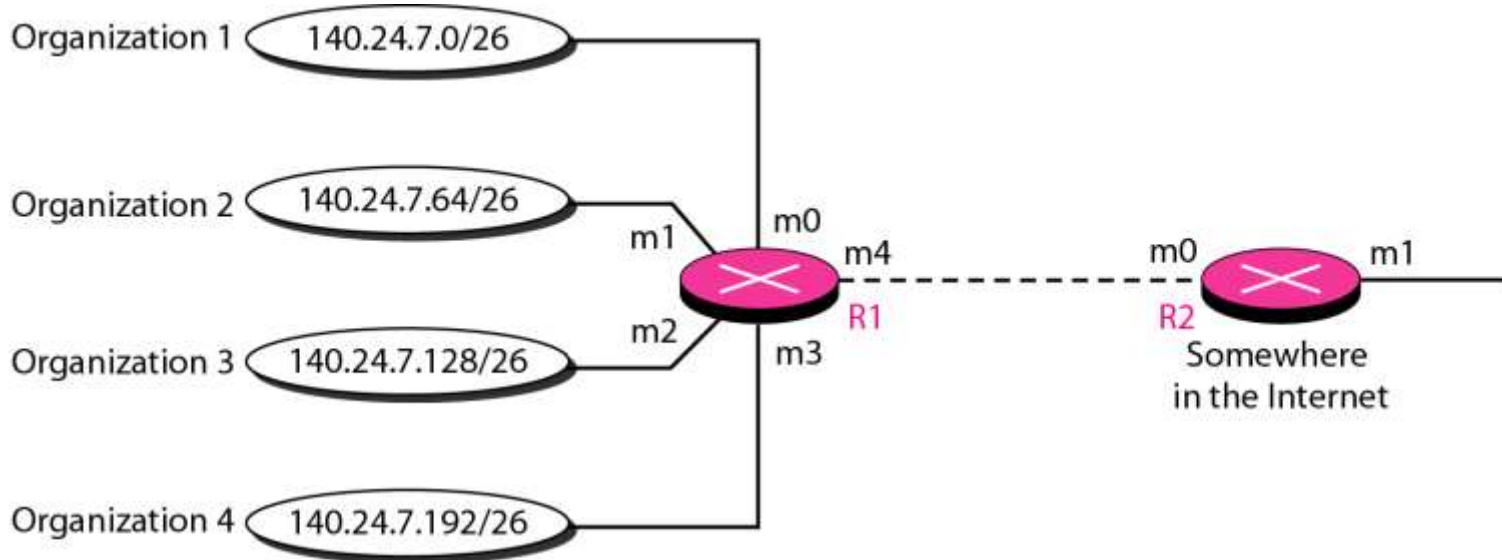
Example-4

Show the forwarding process if a packet arrives at R1 in Figure-6 with the destination address 18.24.32.78.

Solution

This time all masks are applied, one by one, to the destination address, but no matching network address is found. When it reaches the end of the table, the module gives the next-hop address 180.70.65.200 and interface number m2 to ARP. This is probably an outgoing package that needs to be sent, via the default router, to someplace else in the Internet.

Figure-7 *Address aggregation*



| Mask | Network address | Next-hop address | Interface |
|------|-----------------|------------------|-----------|
| /26 | 140.24.7.0 | ----- | m0 |
| /26 | 140.24.7.64 | ----- | m1 |
| /26 | 140.24.7.128 | ----- | m2 |
| /26 | 140.24.7.192 | ----- | m3 |
| /0 | 0.0.0.0 | Default | m4 |

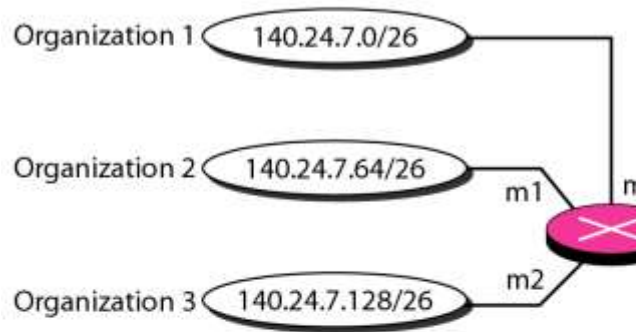
Routing table for R1

| Mask | Network address | Next-hop address | Interface |
|------|-----------------|------------------|-----------|
| /24 | 140.24.7.0 | ----- | m0 |
| /0 | 0.0.0.0 | Default | m1 |

Routing table for R2

Figure-8

Longest mask matching

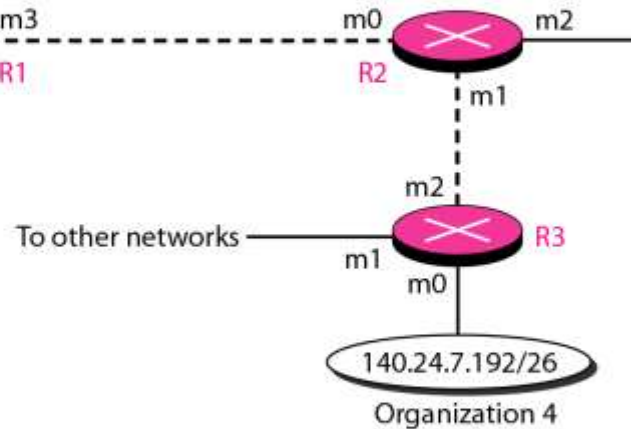


| Mask | Network address | Next-hop address | Interface |
|------|-----------------|------------------|-----------|
| /26 | 140.24.7.0 | ----- | m0 |
| /26 | 140.24.7.64 | ----- | m1 |
| /26 | 140.24.7.128 | ----- | m2 |
| /0 | 0.0.0.0 | Default | m3 |

Routing table for R1

Routing table for R2

| Mask | Network address | Next-hop address | Interface |
|------|-----------------|------------------|-----------|
| /26 | 140.24.7.192 | ----- | m1 |
| /24 | 140.24.7.0 | ----- | m0 |
| /?? | ???????? | ?????????? | m1 |
| /0 | 0.0.0.0 | Default | m2 |



| Mask | Network address | Next-hop address | Interface |
|------|-----------------|------------------|-----------|
| /26 | 140.24.7.192 | ----- | m0 |
| /?? | ???????? | ?????????? | m1 |
| /0 | 0.0.0.0 | Default | m2 |

Routing table for R3

Example-5

As an example of hierarchical routing, let us consider Figure-9. A regional ISP is granted 16,384 addresses starting from 120.14.64.0. The regional ISP has decided to divide this block into four subblocks, each with 4096 addresses. Three of these subblocks are assigned to three local ISPs; the second subblock is reserved for future use. Note that the mask for each block is /20 because the original block with mask /18 is divided into 4 blocks.

The first local ISP has divided its assigned subblock into 8 smaller blocks and assigned each to a small ISP. Each small ISP provides services to 128 households, each using four addresses.

Example-5 (continued)

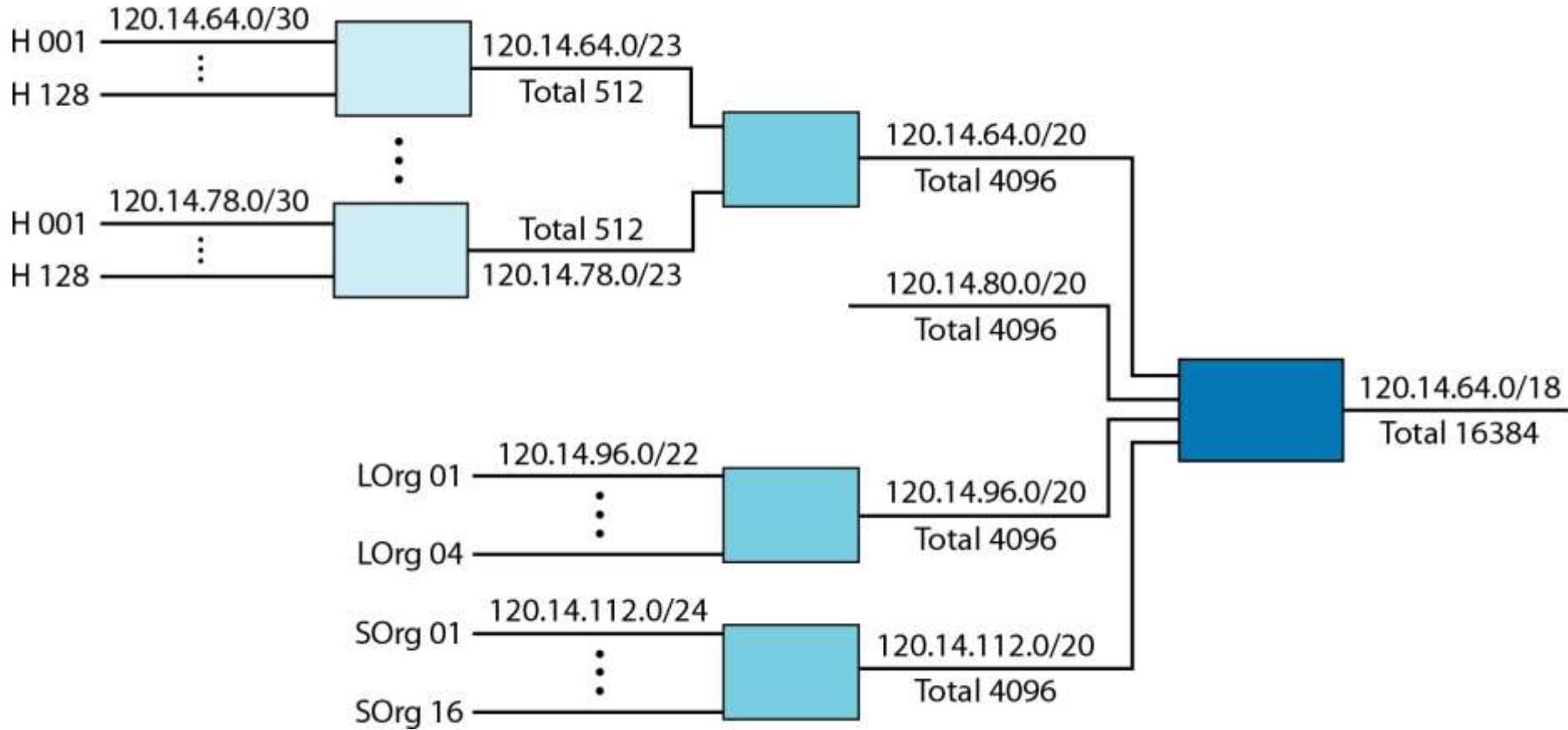
The second local ISP has divided its block into 4 blocks and has assigned the addresses to four large organizations.

The third local ISP has divided its block into 16 blocks and assigned each block to a small organization. Each small organization has 256 addresses, and the mask is /24.

There is a sense of hierarchy in this configuration. All routers in the Internet send a packet with destination address 120.14.64.0 to 120.14.127.255 to the regional ISP.

Figure-9

Hierarchical routing with ISPs



References

- 1. Computer Networks, A. S. Tenenbaum, D. J. Wetheral, Pearson India.***
 - 2. Data Communications and Networking, B.A. Forouzan, Tata McGraw Hill Education Private Limited.***
 - 3. Data and Computer Communications, William Stallings, Pearson-Prentice Hall.***
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